
MFS605/EE605
Systems for Factory Information and Control

Lecture 4

Fall 2005

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UK Center for Manufacturing

1

• **References:**

- **Modeling and Analysis of Mfg. Systems, Askin and Standridge (John Wiley, 1993)**
- **Factory Physics, by Hopp and Spearman (McGraw Hill 1996)**
- **Simulation Modeling and Analysis, 2nd ed. Law and Kelton, 1991 (McGraw Hill)**
- **Probability and Random Processes for Electrical Engineering, 2nd edition, A. Leon-Garcia, 1994 (Addison-Wesley)**
- **Nelson, “Stochastic Modeling, Analysis, and Simulation”, 1995**
- **Hillier & Lieberman, “Intro to Oper. Research, 6th edition”, 1995**

2

Randomness...(Review)

Sources of Variability:

- Natural variation
 - operator difference, tool wear, variation of material
- Breakdown/repair and other unexpected delays
- Setups and other irregular but expected delays
- Quality problems...

Problem: Variability in unbuffered lines lead to blocking and starvation, and thus loss of capacity

3

Random Processing Times without buffers

- Effect of Random Processing Times in balanced, *unbuffered* lines with no breakdowns (Askin/Standridge)
- CV is Coefficient of Variation = (std.dev.) / mean
- → Reducing CV reduces our losses

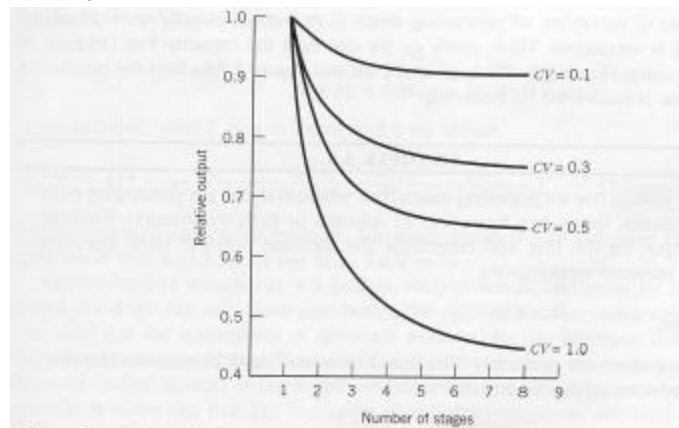


Figure 3.6 Effect of random processing time in balanced, unbuffered lines.

4

Recovery of lost output from buffering (review)

Buffers help reduce starvation and blocking due to time variations and disruptions.

→ Recovery of lost output, but at cost of adding WIP

→ Higher the variation, the more WIP required to recover

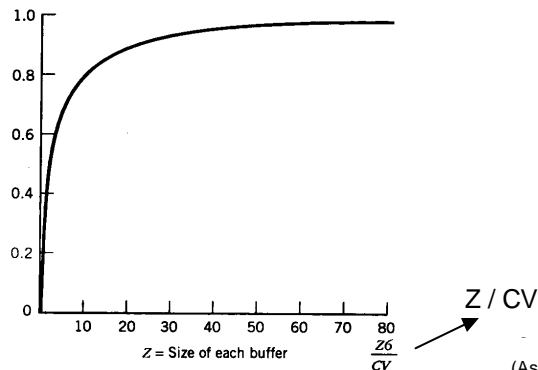


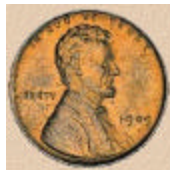
Figure 3.7 Proportion of lost output recovered by buffering in balanced lines.

(Askin and Standridge)

5

Example: Penny Fab 1

(from Hopp and Spearman)



Produce giant novelty pennies.

- Process 1: Punch press cuts penny blanks
- Process 2: Stamping of images
- Process 3: Rim put on penny
- Process 4: Cleaning and deburring

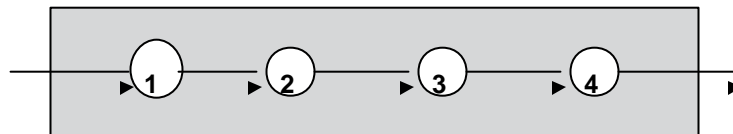
Suppose two hours per machine (“perfectly balanced”)

24 hours /day

$r_b = ?$

Raw Process Time, $T_0 = ?$

Critical WIP = ?



6

Best Case

Best case (zero variability) performance:

- Minimum flow time for a given WIP given by:

$$FT_{\text{best}} \begin{cases} = T_0 & \text{if } w < \text{critical} & \text{(no waiting)} \\ = w / r_b & \text{else} \end{cases}$$

Maximum Throughput rate for a given WIP given by:

$$\text{rate}_{\text{best}} \begin{cases} = w / T_0 & \text{if } w < \text{critical} \\ = r_b & \text{else} \end{cases}$$

This is the case for PennyFab 1
(perfectly balanced, no variability)

7

Worst Case (impractical)

Worst Case (impractical): Suppose average processing time per station was same, but actually all the time was spent with the first job.

Note: this is still not random, but we have picked the worst case possible

Example: 1st job = 8hrs; 2nd, 3rd, 4th take 0 hours.

Then: avg.=2 still, but 4th job must wait 8 hours at each station!

For given WIP w , define:

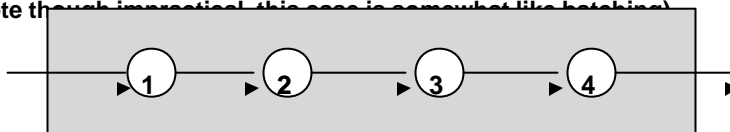
$$\begin{aligned} FT_{\text{worst}} &= w T_0 && \text{(waiting on all other WIP)} \\ \text{rate}_{\text{worst}} &= 1/T_0 \end{aligned}$$

Example: For Penny Fab 1, we have to wait 8 hours each station

For $w = 4$, Flow time = $8 + 8 + 8 + 8 = 32$ hours = $4 * 8$

Throughput rate = $4/32 = 1/8$ jobs / hour

(Note though impractical, this case is somewhat like batching)



8

Practical Worst Case

- Real systems rarely operate like the best or worst case.
- To consider between the best and worst, we invoke probability.
- The derivation of this “practical worst case” will be done later. For now, we present the basic idea:
 - All stations perfectly balanced
 - All stations are single machines
 - Process times are random. In particular, assume that our system is “memoryless” – so the remaining time for a job at a station is independent of the time that the job has already been there.
- See graphs next page – derivations to be done with queueing material

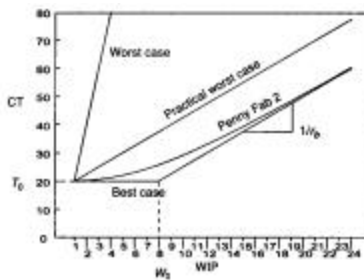
9

WIP relationships

(Hopp and Spearman, 1996)

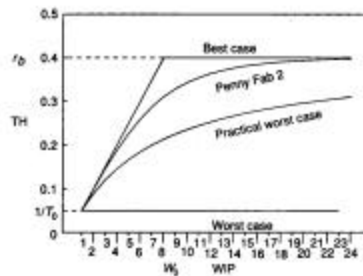
- Flow time vs. WIP

FIGURE 7.9
Cycle time versus WIP
in Penny Fab Two



- rate vs. WIP

FIGURE 7.10
Throughput versus
WIP in Penny
Fab Two



10

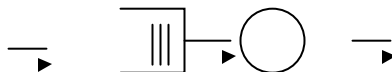
Queueing models

- Queueing Models
 - Probabilistic -- primarily steady state averages
 - Advantages:
 - Analytical solution to problems with randomness & uncertainty
 - --> fast solution (*if solution exists or is known!*)
 - Disadvantages:
 - Requires simplifying assumptions
 - Best for smaller models
 - solutions not always possible
 - most suited for steady state analysis

11

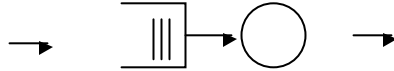
Queueing Theory: Basic Terms

- Queueing Theory: Study of lines, waiting in lines



- Buffer or Queue
- Server -- service rate = m
 - machine
 - repair
 - (sometimes even parts)
- Customers --- arrival rate = l
 - parts waiting for machining
 - machines waiting for repair
 - ...

12



- **Issues: What is distribution of arrivals**
 - What is distribution of service time
 - What is policy for selecting next customer?

- **Queuing theory answers**
 - expected number of customers in system
 - expected number of customers in queue
 - avg. customer waiting time in system
 - avg customer waiting time in queue
 - avg. utilization of server

13

Queuing Theory

- In this course, we focus *mostly* on simplest queue: **M/M/1**
 - **M** stands for *Markov* or *Memoryless*
 - **Notation: Arrival Process / Service Process / # of servers**

 - **M/M/1 -->**
 - Poisson arrivals: (exponential interarrival times)
time distribution = $\lambda e^{-\lambda t}$ for $t \geq 0$
 λ = mean arrival rate (expected # arrivals / time)
 - Exponential service times
time distribution = $\mu e^{-\mu t}$ for $t \geq 0$
 μ = mean service rate (expected # of customers completing service per unit time).

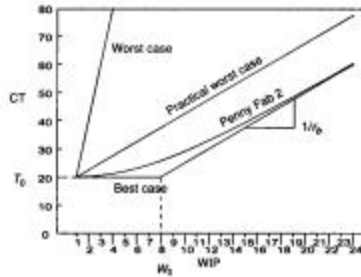
14

WIP relationships (review)

(Hopp and Spearman, 1996)

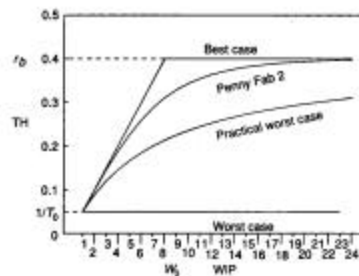
- Flow time vs. WIP

FIGURE 7.9
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in Penny Fab Two



- rate vs. WIP

FIGURE 7.10
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15

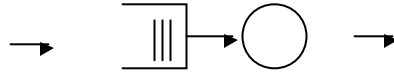
“Practical Worst Case”

- The “practical worst case” in previous diagrams corresponds to a balanced line of single machine stations with exponential process times.
 - (even worse can happen under bad batching)
- “Exponential distribution” of times is “Markov” or “Memoryless”.
 - This extreme randomness arises from:
 - “random” arrivals (time to next is not given by time since prev.)
 - Highly variable service times – examples:
 - job shop or repair station
 - Setups
- We will study models of behavior for this “practical worst case”
 - → Queueing models
 - Provides analytical relationships of average waits, average lengths of queues, etc.
 - Results on only restricted classes of systems

16

Queuing Theory: Basic Terms

- Queuing Theory: Study of lines, waiting in lines



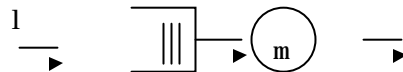
- Buffer or Queue
- “Server” -- service rate = m
- “Customers” --- arrival rate = l

In this course, we focus *mostly* on simplest queue: M/M/1

- (Notation: Arrival Process / Service Process / # of servers)
- M is for “markov” or “memoryless”
 - → Poisson arrivals (exponential inter-arrivals)
 - → Exponential service times
- “First Come First Served” (FIFO) queue discipline

17

Utilization Assumption



- We assume $l/m < 1$.
 - Why??

$r = l/m$ is called *utilization*

18

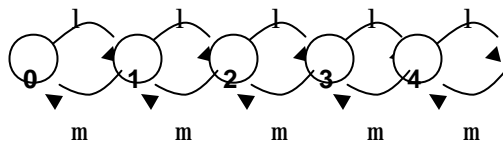
Exponential Assumption

- Exponential distributions give simplest queuing analysis
- Is exponential distribution reasonable?
 - For service time distribution:*
 - inappropriate if service same over most parts (as if same parts)
 - appropriate if variety of customers each requiring different operations
 - For Interarrival time distributions:*
 - Appropriate if customers appear in very random manner
 - Inappropriate if customers typically appear in groups, or if customers may postpone arrival based on length of queue

19

Birth-Death Model

- Quantities of interest in queuing theory:
 - P_n : probability of having n parts in system
 - W : avg. wait in system
 - L : avg. number of customers in system
- Consider steady-state averages*
Assume unbounded queue size

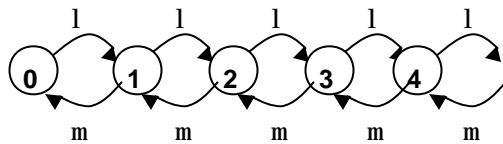


Consider steady-state averages: --> Freq. of entry to state same as frequency of departing from state.

20

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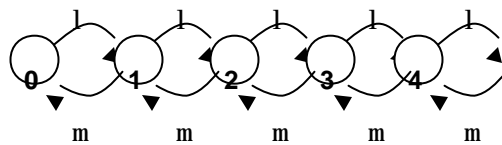


Consider steady-state averages: \rightarrow Freq. of entry to state same as frequency of departing from state.

21

Deriving P_n from birth-death model

- Balance Equations: entry freq. = exit freq. for each state.



- Insert derivations

22

Results

- $P_0 = 1 - l/m = 1 - r$
- $P_n = (l/m)^n (1 - l/m) = r^n (1 - r) = r^n P_0 \quad (\text{for } n > 0)$
- **Example 1:** $l = 5$ parts/hr, $m = 10$ parts / hr.
- **Example 2:** $l = 8$ parts/hr, $m = 10$ parts / hr.

23

Other questions:

- What is probability that the machine is idle?
- What is the probability that there are more than 3 customers in the system?

24

Avg. # of customers

- L : = avg. number of customers *in the system*

$$L = E[n] = \sum_{n=0}^{\infty} nP_n$$

$$L = \lambda / (\mu - \lambda) = r / (1 - r)$$

25

Little's Law

- W = mean time spent in system

- Little's law:

$$L = \lambda_{\text{eff}} W$$

mean number of customers in system =
effective arrival rate x mean wait time

mean # in "system" = arrival rate x mean wait in "system"

mean # in queue = arrival rate x mean wait in queue

Why the "effective arrival rate"?

26

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- Nice graphical proof of Little's Law found in Leon-Garcia, 2nd edition, pages 502-503.

Mean Wait Time in System

- W = mean wait time in system

$$L = \frac{\lambda}{(m-1)} = \frac{\lambda}{(1-r)}$$

- Little's law: $L = \lambda W$

- $W = L / \lambda = 1 / (m-1)$

- Examples:

-
- Example: We have demand rate of 1 part per hour. We want average lead time per part under 2 hours. How fast a workstation do we need?

29

Summary of results so far

M/M/1 queue:

- exponential service times (service rate = m)
- Poisson arrivals (arrival rate = l)
- 1 server, infinite buffer capacity

- $P_0 = 1 - l/m = 1 - r$
- $P_n = (l/m)^n (1 - l/m) = r^n (1 - r) = r^n P_0$ (for $n > 0$)

- **L:** = mean number of customers *in the system*

$$L = E[n] = \sum_{n=0}^{\infty} n P_n = \frac{l}{(m-l)} = \frac{r}{1-r}$$

- **W:** = mean wait time in the system

$$W = \frac{L}{l} = \frac{1}{(m-l)} = \frac{1}{m(1-r)}$$

30

Waiting time in Queue and # in Queue

- W_Q : mean waiting time in just the queue
 = mean time in system - mean time in service
 $W_Q = W - 1/m$

- L_Q : mean number of customers in the queue

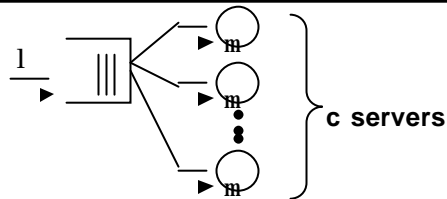
- Use Little's Law:

$$L_Q = \lambda W_Q$$

Examples:

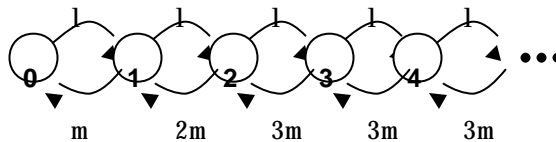
31

Multiple servers: M/M/c



- Utilization: $r = \frac{\lambda}{cm} < 1$

- Basic intuition from birth-death graph (for $c=3$):



32

Basic results for M/M/c

- (table 11.1 from Askin Standridge)

	M/M/1	M/M/c
P_0	$1 - \rho$	$\left[\frac{(c\rho)^c}{c!(1-\rho)} + \sum_{n=0}^{c-1} \frac{(c\rho)^n}{n!} \right]^{-1}$
L_q	$\frac{\rho^2}{1-\rho}$	$\frac{\rho(c\rho)^c P_0}{c!(1-\rho)^2}$
L	$\frac{\rho}{1-\rho}$	$L_q + \frac{\lambda}{\mu}$
W_q	$\frac{\rho}{\mu(1-\rho)}$	$\frac{(c\rho)^c P_0}{c!c\mu(1-\rho)^2}$
W	$\frac{1}{\mu(1-\rho)}$	$W_q + \mu^{-1}$

33

Example of M/M/c

- Suppose $c=2$, $m = 5$, $l = 8$

$$r = \frac{l}{cm} = \frac{8}{2 \cdot 5} = \frac{8}{10}$$

$$P_0 = \left(\frac{(cr)^c}{c!(1-r)} + \sum_{n=0}^{c-1} \frac{(cr)^n}{n!} \right)^{-1} = \left(\frac{(1.6)^2}{2 \cdot (2)} + \left(1 + \frac{1.6}{1} \right) \right)^{-1} = (9)^{-1} = 0.1111$$

$$L_q = \frac{r(cr)^c P_0}{c!(1-r)^2} = \frac{0.8(1.6)^2(0.1111)}{2 \cdot (2)^2} = 2.84$$

$$L = L_q + \frac{l}{m} = 2.84 + 1.6 = 4.44$$

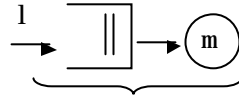
$$W = \frac{L}{l} = \frac{4.44}{8} = 0.555$$

$$W_q = \frac{L_q}{l} = \frac{2.84}{8} = 0.355$$

34

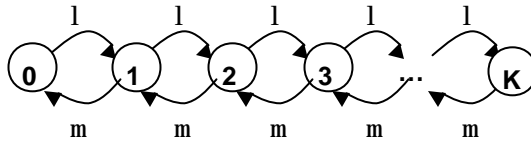
M/M/1/K

- Finite capacity queues: State probabilities and # in system



System capacity of K

- Birth-death model



$$P_i = r^i P_0$$

$$1 = \sum_{i=0}^K P_i = \sum_{i=0}^K r^i P_0 \longrightarrow P_0 = \frac{1-r}{1-r^{K+1}}$$

35

$$L = E[n] = \sum_{n=0}^K n P_n = \begin{cases} \frac{r}{1-r} - \frac{(K+1)r^{K+1}}{1-r^{K+1}} & \text{for } r \neq 1 \\ \frac{K}{2} & \text{for } r = 1 \end{cases}$$

- M/M/1/K queues: mean wait in system:**
 - Note some customers get turned away.
 - For Little's Law, we need **effective arrival rate**
 - offered load is measure of demand on system: λ/μ
 - carried load is actual demand met by system: λ_a/μ

$$I_a = \text{effective arrival rate} = I(1 - P_K)$$

$$W = \frac{L}{I_a}$$

$$W_Q = W - \frac{1}{m}$$

$$L_Q = W_Q I_a \quad \text{note it can be shown this is } L_Q = L - (1 - P_0)$$

36

M/M/1/K example

- Suppose $K=6$, $m = 1/8$, $l = 1/10$

$P_i = ?$

M/M/1/6	M/M/1/infinity
P0 = .25307	P0=.2
P1=.20246	P1=.16
P2=.16197	P2=.128
P3=.12957	P3=.1024
P4=.10366	P4=.08192
P5=.08293	P5=.065536
P6=.06634	P6=.0524288
Total P0 to P6: 1.0	Total P0 to P6: .790284

- Key point: the limited size system turns away parts sometimes since the M/M/1/infinity system has probability $(1-.79)= 21\%$ of being in state beyond 6.

37

M/G/1 Queues

$$V(N) = E(N) + \frac{\lambda^2 E(S^2)}{2(1-\rho)} + \frac{\lambda^4 [E(S^2)]^2}{4(1-\rho)^2} \quad (11.9)$$

- (from page 367, Askin and Standridge) – FCFS (FIFO)

$$E(T) = E(S) + \frac{\lambda E(S^2)}{2(1-\rho)} \quad (11.6)$$

$$V(T) = V(S) + \frac{\lambda E(S^3)}{3(1-\rho)} + \frac{\lambda^2 [E(S^2)]^2}{4(1-\rho)^2} \quad (11.7)$$

$$E(N) = \rho + \frac{\lambda^2 E(S^2)}{2(1-\rho)} \quad (11.8)$$

$$V(N) = E(N) + \lambda^2 V(S) + \frac{\lambda^3 E(S^3)}{3(1-\rho)} + \frac{\lambda^4 [E(S^2)]^2}{4(1-\rho)^2} \quad (11.9)$$

- S is for service time (so $E(S) = 1/\mu$)
- T is for throughput time (so $E(T) = W$)
- N is for number of jobs (so $E(N) = L$)

Note reducing $E(S^2)$ reduces $E(T)$ and $E(N)$

Note that $E(T)$ and $E(N)$ depend only of variance and mean of service time distribution, not its distributional form.

38

M/G/1 Queues

- (from page 367, Askin and Standridge)

$$E(T) = E(S) + \frac{\lambda E(S^2)}{2(1-\rho)} \quad (11.6)$$

$$V(T) = V(S) + \frac{\lambda E(S^3)}{3(1-\rho)} + \frac{\lambda^2 [E(S^2)]^2}{4(1-\rho)^2} \quad (11.7)$$

$$E(N) = \rho + \frac{\lambda^2 E(S^2)}{2(1-\rho)} \quad (11.8)$$

Example from text: $\lambda = 10/\text{week}$. All distributions with same mean
Key points: the exponential distribution gives the worst case out
 the distributions with comparable means. It is thus useful as
 a practical worst case (maximum randomness) estimate.

Table 11.2 Performance Comparisons for M/G/1 Example

Distribution	$E(S)$	$V(S)$	$E(T)$	$V(T)$	$E(N)$	$V(N)$
Exponential	0.0833	0.00694	0.500	0.250	5.00	30.0
Uniform	0.0833	0.0023148	0.361	0.103	3.61	13.9
Gamma	0.0833	0.000694	0.313	0.0685	3.13	9.97
Deterministic	0.0833	0.0	0.292	0.0550	2.92	8.41

39

Review:

- M/M/1 queue:**

$$L = \lambda / (m - \lambda) = r / (1 - r)$$

$$W = L / \lambda = 1 / (m - \lambda)$$

$$W_Q = W - 1/m$$

$$L_Q = \lambda W_Q$$

- M/M/1/K**

$$L =$$

$$\lambda_a = \lambda (1 - P_K)$$

$$W = L / \lambda_a$$

$$W_Q = W - 1/m$$

$$L_Q = \lambda_a W_Q$$

40

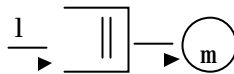
**Effect of many small (many queue) vs.
~~single machine~~**

41

**Effect of many small (multi-queue) vs.
~~multi-server (same queue)~~**

42

Summary



- **M/M/1 Queues**
- **M/M/c Queues** ---- **c servers**
- **M/G/1 Queues** ---- **generalized service distributions**
- **M/M/1/K Queues** ---- **Limited capacity in system**

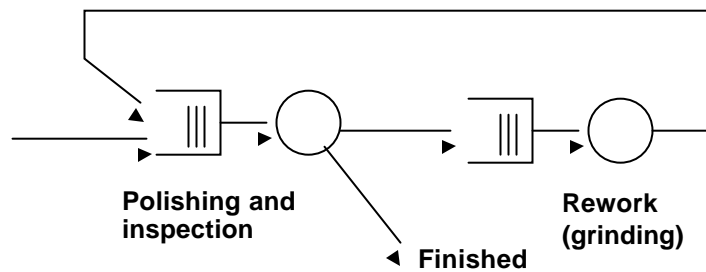
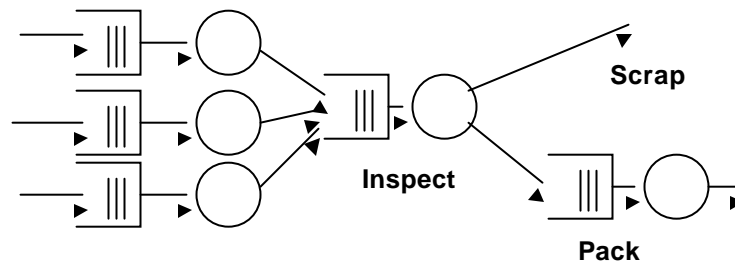
Systems with multiple queues

- Open networks vs. closed networks

- open Jackson networks:
 - open network
 - Poisson external arrivals
 - exponential service times (possibly $c > 1$)
 - unlimited buffers at each node
 - probabilistic routing
 - First come – First Serve
 - one type of customer

45

Examples



46

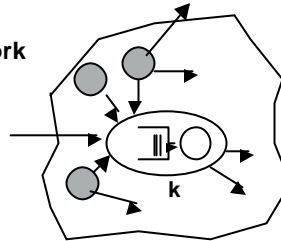
Open Jackson Networks

- Why Jackson networks?
 - Probability of system state can be calculated from looking at the individual queues in the system.

Under steady state conditions, each facility in a Jackson network behaves as if it were an independent M/M/s queuing system with effective arrival rate:

$$\bullet \lambda_k = a_k + \sum_{i=1..m} \lambda_i p_{ik}$$

- p_{ik} probability of routing from i to k
- a_k arrival rate in from outside
- λ_k arrival rate at node k
- m number of nodes in network



47

Jackson Network Solution Procedure

1. Solve for arrival rates using

$$\lambda_k = a_k + \sum_{i=1..m} \lambda_i p_{ik}$$

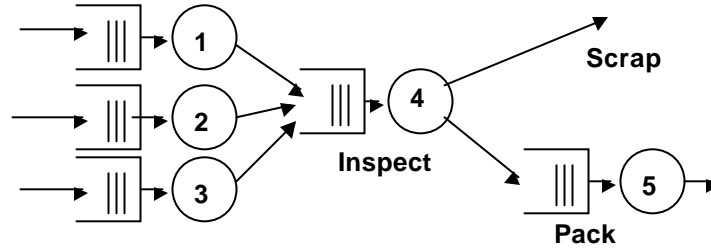
1. Analyze each workstation as a M/M/1 or M/M/c queue

Use standard table of L, W, L_Q , W_Q formulas

2. Combine results across workstations to get performance measures of the entire system.

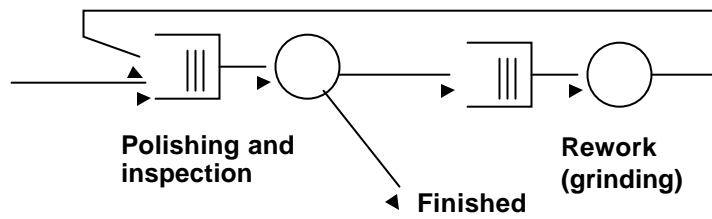
48

Queueing Network example 1



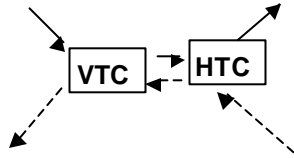
49

Queueing Network example 2



50

Queuing network example 3



	<i>Part A</i>	<i>Part B</i>
Path	VTC-HTC-exit	HTC-VTC-exit
Demand	Poisson 30/day	Poisson 60/day
VTC time (min.)	uniform(8,12)	uniform(8,12)
HTC time(min.)	uniform(10,14)	uniform(6,10)

Extensions

Networks with general distribution service times.

